Polynomial Functors and Shannon Entropy

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Outline

1 Introduction

- Why am I here?
- Working in the Poly ecosystem
- Plan of the talk

2 Background on Poly

3 Distributive functors and entropy

4 Conclusion

Why am I here?

We're here to learn from each other. But what is learning?

- Somehow out of all the information out there, some of it *sticks*.
- We develop frameworks by which to *store* information.
- I'm interested in how intelligence and learning function.
- So I study how knowledge is stored and transferred in databases and...
- ...how dynamical systems interact to adapt and learn (e.g. in DNNs).

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Entropy has been put forward as an approach to intelligence and learning.

- Life can be understood as a *dissipative system*, spraying entropy.
- It does so while packing negentropy—organization—into itself.
- Polani's empowerment and Freer's causal entropic forces...
- ...are entropy-based approaches to intelligent behavior.

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- Life can be understood as a *dissipative system*, spraying entropy.
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- Polani's empowerment and Freer's causal entropic forces...
- ...are entropy-based approaches to intelligent behavior.
- I only seem to understand things when they're written categorically.
 - I've been trying to understand "what entropy really is."
 - The Baez-Fritz-Leinster conception of entropy is great,...
 - ...but I want to connect it in with dynamical systems or databases.

The overwhelming abundance of Poly

In January 2020 I fell in love with a category called Poly.

- Its applications subsume everything I'd done with categ'l databases...
- ...and everything I'd done with interacting dynamical systems.
- It's used in functional programming, type theory, higher cat'y theory.

The overwhelming abundance of Poly

In January 2020 I fell in love with a category called **Poly**.

- Its applications subsume everything I'd done with categ'l databases...
- ...and everything I'd done with interacting dynamical systems.
- It's used in functional programming, type theory, higher cat'y theory.
- But it's not just very applicable, it's also very highly-structured.
 - Coproducts and products that agree with usual polynomial arithmetic;
 - All limits and colimits;
 - At least three orthogonal factorization systems;
 - A symmetric monoidal structure \otimes distributing over +;
 - A cartesian closure q^p and monoidal closure [p, q] for \otimes ;
 - Another nonsymmetric monoidal structure <> that's duoidal with <>;
 - A left (Meyers?)

-coclosure $\begin{bmatrix} \\ \end{bmatrix}$, meaning $Poly(p, q \triangleleft r) \cong Poly(\begin{bmatrix} r \\ p \end{bmatrix}, q)$;
 - An indexed right <-coclosure, i.e. $\operatorname{Poly}(p, q \triangleleft r) \cong \sum_{f \colon p(1) \rightarrow q(1)} \operatorname{Poly}(p \frown q, r);$
 - An indexed right \otimes -coclosure (Niu?), i.e. $\operatorname{Poly}(p, q \otimes r) \cong \sum_{f: p(1) \to q(1)} \operatorname{Poly}(p \nearrow q, r);$
 - At least ten monoidal structures in total;
 - ⊲-monoids generalize Σ-free operads;
 - <-comonoids are exactly categories; bicomodules are data migrations.

See "A reference for categorical structures on Poly", arXiv: 2202.005342/15

Entropy in terms of Poly

Today I'll tell you how entropy looks from the **Poly** point of view.

- I'll show how to think of objects in **Poly** as empirical distributions.
- I'll show that there are distributive monoidal functors

$$\mathsf{Poly}^{\mathsf{Cart}} \xrightarrow{p \mapsto \dot{p}y} \mathsf{Poly} \xrightarrow{p \mapsto (p(1), \Gamma(p))} \mathsf{Set} \times \mathsf{Set}^{\mathsf{op}}$$

sending p ∈ Poly^{Cart} to an invariant h(p) := (A, B) ∈ Set × Set^{op}.
The Shannon entropy can then be extracted: H(p) = log(A/ ^A√B).
Properties of entropy follow from the distributive monoidality of h.

Plan

The plan for the rest of the time is as follows:

- Give background on polynomial functors.
- Explain $h: \mathbf{Poly}^{\mathbf{Cart}} \to \mathbf{Set} \times \mathbf{Set}^{\mathsf{op}}$ and its relation to entropy.
- Talk about generalizations and future work.
- Conclude.

Outline

1 Introduction

2 Background on Poly

- The category Poly
- Distributive monoidal structure
- Other theoretical aspects

3 Distributive functors and entropy

4 Conclusion

What is a polynomial?

Here's how we'll think about polynomial functors.



What is a polynomial?

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You can think of the bundle as a empirical distribution:

- The first outcome was drawn twice; the next three once; the rest never.
- It corresponds to the distribution $(\frac{2}{5}, \frac{1}{5}, \frac{1}{5}, \frac{1}{5}, 0, 0)$.

What is a morphism of polynomials?

Let
$$p \coloneqq y^3 + 2y$$
 and $q \coloneqq y^4 + y^3 + y + 1$



A morphism $p \xrightarrow{\varphi} q$ sends *p*-outcomes to *q*-outcomes, interpreting draws:



It is Cartesian iff each map on draws is a bijection.

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It is *Cartesian* iff each map on draws is a bijection. (**Poly**^{Cart} $(p,q) \cong 6$.)

Fundamental invariants

 $2y^3 + y^2 + y =$

We will be interested in two fundamental invariants of a polynomial.

- From the bundle POV, these would be *base* and *global sections*.
- So if p is represented by $E \to B$, these are B and $\mathbf{Set}_{/B}(B, E)$.

In terms of polynomials these are

 $p(1) \cong \operatorname{\mathsf{Poly}}(y,p)$ and $\Gamma(p) \coloneqq \operatorname{\mathsf{Poly}}(p,y).$

• E.g. for the following bundle these are $p(1) \cong 4$ and $\Gamma(p) \cong 18$.



has 4 base elements and 18 global sections

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These are functorial in opposite directions:

$$\textbf{Poly} \xrightarrow{p \mapsto (p(1), \Gamma(p))} \textbf{Set} \times \textbf{Set}^{op}$$

This functor has a right adjoint, but we won't use that.

The distributive monoidal structure (Poly, $0, +, y, \otimes$)

The category **Poly** is distributive monoidal.

- The usual sum of two polynomials is their coproduct; 0 is initial.
- There is another operation \otimes called *Dirichlet product*. Formula:

$$p \otimes q \coloneqq \sum_{(i,j) \in p(1) imes q(1)} y^{p[i] imes q[j]}$$

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These clearly distribute: p ⊗ (q₁ + q₂) ≅ (p ⊗ q₁) + (p ⊗ q₂).
Soon we'll see how the fundamental invariants respect these oper'ns.

Derivatives and the bifibration

The derivative of a polynomial functor is another polynomial functor.

• Write \dot{p} for the derivative with respect to y; consider also $\dot{p}y$.

$$\dot{p} = \sum_{i \in p(1)} \sum_{d \in p[i]} y^{p[i] \setminus \{d\}}$$
 and $\dot{p}y \cong \sum_{i \in p(1)} p[i] y^{p[i]}$

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Neither of these is functorial in **Poly**, but *they are functorial* in **Poly**^{Cart}.

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The last theory we'll need is the bifibration $\textbf{Poly} \rightarrow \textbf{Set}.$

- The functor $p \mapsto p(1)$ is a monoidal *-bifibration. All we need:
- For any $f: A \to B$, we have $\mathbf{Poly}_A(f^*q, p) \cong \mathbf{Poly}_B(q, f_*p)$.

$$f^*q \coloneqq \sum_{a \in A} y^{q[fa]}$$
 and $f_*p \coloneqq \sum_{b \in B} y^{\sum_{a \mapsto b} p[a]}$

As we lump outcomes together, we add up the draws.

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3 Distributive functors and entropy

- Distributive functor $\mathbf{Poly}^{\mathbf{Cart}} \rightarrow \mathbf{Set} \times \mathbf{Set}^{\mathsf{op}}$
- Entropy and entropy density

4 Conclusion

$p\mapsto \dot{p}y$ is distributive

We're now ready to get to work on how all this relates to entropy.

- The approach is to extract two invariant sets from any polynomial.
- This process is "good" in that it is a distributive monoidal functor.
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So first, we want to see that $p\mapsto \dot{p}y$ is distributive monoidal.

• The derivative is linear, $(p+q) = \dot{p} + \dot{q}$, and so is $p \mapsto py$.

• So $p \mapsto \dot{p}y$ preserves coproducts +. It also preserves \otimes :

$$\begin{aligned} (\dot{p}y) \otimes (\dot{q}y) &\cong \sum_{i \in p(1)} p[i]y^{p[i]} \otimes \sum_{j \in q(1)} q[j]y^{q[j]} \\ &\cong \sum_{\substack{(i,j) \in p(1) \times q(1) \\ G &: G \\ &\cong (p \otimes q)y}} p[i] \times q[j]y^{p[i] \times q[j]} \end{aligned}$$

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$$\cong \sum_{\substack{(i,j) \in p(1) \times q(1)\\ \ominus}} p[i] \times q[j]y^{p[i] \times q[j]}$$
$$\cong (p \otimes q)y$$

So (**Poly**^{Cart}, 0, +, y, \otimes) $\xrightarrow{p \mapsto \dot{p}y}$ (**Poly**, 0, +, y, \otimes) is distributive monoidal.

Fundamental invariants are distributive

The fundamental invariants $p \mapsto (p(1), \Gamma(p))$ are also distributive

$$(\mathbf{Poly},0,+,y,\otimes) \xrightarrow{\boldsymbol{p} \mapsto (\boldsymbol{p}(1), \boldsymbol{\Gamma}(\boldsymbol{p}))} (\mathbf{Set} \times \mathbf{Set}^{\mathsf{op}}, (0,1),+, (1,1), \otimes)$$

But what exactly is all this structure on $\mathbf{Set} \times \mathbf{Set}^{\mathrm{op}}$?

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But what exactly is all this structure on $\mathbf{Set} \times \mathbf{Set}^{op}$?

 Set × Set^{op} has coproducts: (A₁, B₁)+(A₂, B₂) ≃ (A₁+A₂, B₁×B₂).
 It has another symmetric monoidal structure with unit (1, 1): (A₁, B₁) ⊗ (A₂, B₂) := (A₁ × A₂, B₁^{A₂} × B₂^{A₁})

• And these distribute "because" $B^{A_1+A_2}(B_1B_2)^A \cong (B^{A_1}B_1^A)(B^{A_2}B_2^A).$

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$$(A_1,B_1)\otimes (A_2,B_2)\coloneqq (A_1\times A_2,B_1^{A_2}\times B_2^{A_1})$$

And these distribute "because" $B^{A_1+A_2}(B_1B_2)^A \cong (B^{A_1}B_1^A)(B^{A_2}B_2^A)$. Why do the fundamental invariants (as a pair) preserve + and \otimes ?

- We have $(p+q)(1) \cong p(1) + q(1)$ and $\Gamma(p+q) \cong \Gamma(p) \times \Gamma(q)$.
- This says they preserve +. One also checks they preserve ⊗:

 $(p\otimes q)(1)\cong p(1) imes q(1)$ and $\Gamma(p\otimes q)\cong \Gamma(p)^{q(1)} imes \Gamma(q)^{p(1)}$

Taking stock

Let's denote the composite of our distributive functors by h:



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Theorem

Let p be a polynomial, considered as a probability distribution P, and let H(P) be its Shannon entropy. Then we have

$$H(P) = L(h(p))$$

The categorical partition function and entropy

I'm unfamiliar with the thermo picture. Joint with James Dama:

- One should think of Shannon entropy as an *entropy density*.
- The thermo picture defines a *partition function* Ω for distributions.
- For $p \in \mathbf{Poly}$ with h(p) = (A, B), this would be $\Omega_p := \frac{A^A}{B}$.
- Then the extensive entropy of p is given by $E(p) := \log \Omega(p)$.
- And the Shannon entropy of p is the density H(p) := E(p)/A.

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Let's see a worked example. Consider the bundle for $p := y^4 + 4y^1$:



• We find $\dot{p}y = 4y^4 + 4y$, so $h(p) = (4 + 4, 4^4) = (8, 4^4)$. • So $\Omega_p = \frac{8^8}{4^4}$, Ext've: $E(p) = \log \Omega_p = 16$, Shannon: H(p) = 16/8 = 2.

Consequences of distributivity

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- Suppose you write p as a sum, $p \coloneqq \sum_{a \in A} p_a$.
- This is the same as giving a function $f: p(1) \rightarrow A$.
- Recall that from this we get $f_*p \in \mathbf{Poly}_A$, lumping distributions.

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- **Recall that from this we get** $f_*p \in \mathbf{Poly}_A$, lumping distributions.

It follows from the fact that h preserves sums that:

$$\Omega_p = \Omega_{f_*p} imes \prod_{a \in A} \Omega_{p_a}$$
 and $\mathsf{E}(p) = \mathsf{E}(f_*p) + \sum_{a \in A} \mathsf{E}(p_a)$

The usual "chain rule" for Shannon entropy follows directly from this.

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- There is a $(+, \otimes)$ -preserving functor h: Poly^{Cart} \rightarrow Set \times Set^{op}.
- Theorem: if (A, B) := h(p) then $H(p) = \log(A/\sqrt[A]{B})$.
- So all the entropy-relevant data of p is encapsulated in two sets.

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- Theorem: if (A, B) := h(p) then $H(p) = \log(A/\sqrt[A]{B})$.
- So all the entropy-relevant data of *p* is encapsulated in two sets. Entropy still feels somehow foreign to me.
 - Hopefully, having different categorifications will help clarify it.
 - I still have hope that it will someday bond with the dynamics of Poly.

Thanks! Comments and questions welcome...