

# TYPE THEORY & THEMES IN PHILOSOPHICAL LOGIC

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<https://consequently.org/p/2025/tt-tpl>

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2 MODAL & SUBSTRUCTURAL LOGICS

3 INTENSIONALITY & IDENTITY

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6 FORMAL & APPLIED TYPE THEORY



# 1 INTRODUCTION

2 MODAL & SUBSTRUCTURAL LOGICS

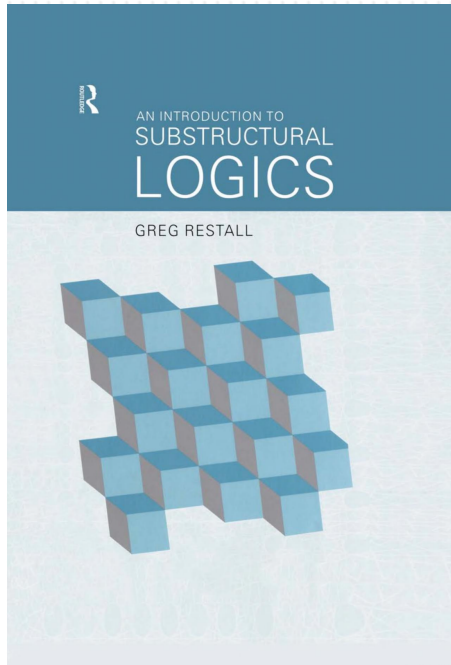
3 INTENSIONALITY & IDENTITY

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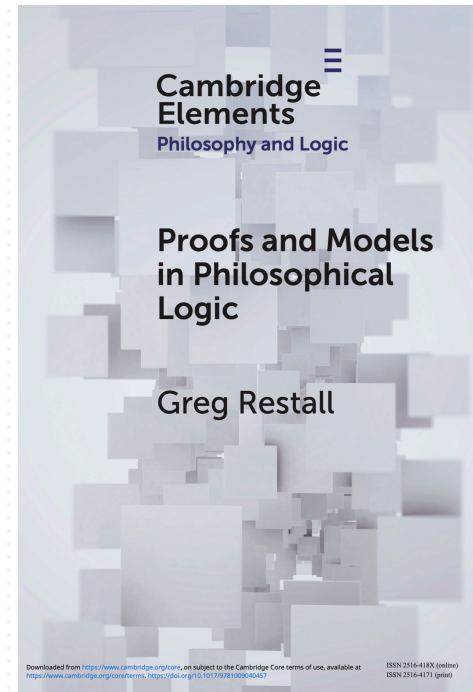
I work in PHILOSOPHICAL LOGIC



2000



2006



2022

I work to understand the connections between different techniques, traditions & approaches in logic & philosophy  
TYPE THEORY is an exciting world I am beginning to explore.

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# MODAL LOGICS — possibility & necessity; reasoning over times, ---

- Massive industry  $\left\{ \begin{array}{l} \cdot \text{"metaphysical necessity"} \\ \cdot \text{Epistemic logics} \\ \cdot \text{Montague style type theory in linguistics} \end{array} \right.$

Central tool: Kripke-style "possible worlds semantics."

Good tools at the level of types, not terms.

(Kripke models represent what follows from what — not why.)

Minority tradition ~ algebras & proof theory for modal logics.

RESIDUATION  
Galois  
Connection

$$\left\{ \begin{array}{c} \frac{\Diamond^{\leftarrow} a \leq b}{a \leq \Box^{\rightarrow} b} \end{array} \right.$$

$\Diamond^{\leftarrow}$  some time in the past

$\Box^{\rightarrow}$  all times in the future

Generalises more naturally to categories & so, to type theoretical interpretation.

# SUBSTRUCTURAL LOGICS — resources, relevance, paradox, syntax

$$\frac{a \otimes b \leq c}{a \leq b \rightarrow c}$$
 } "substructural" since the standard structural rules of contraction, weakening, permutation, & even associativity may be absent.

Kripke models for modal logics generalise to the substructural setting.

$\Box^{\rightarrow}$   $\Diamond^{\leftarrow}$  — unary connective, binary relation  $\Box^{\rightarrow}$  universal forward  $\Diamond^{\leftarrow}$  existential backward  
 $\rightarrow$   $\otimes$  — binary connective, ternary relation  $\rightarrow$  universal forward  $\otimes$  existential backward

These models extend distributive lattices with  $\rightarrow, \otimes$ .

(Algebras, Coherence Spaces & Phase space models give natural non-distributive structures & categories.)

1 INTRODUCTION

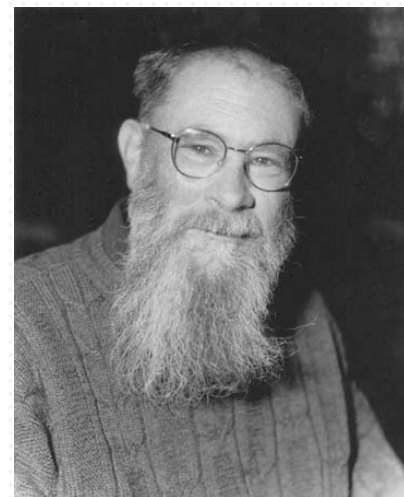
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Identity is utterly simple and unproblematic.  
Everything is identical to itself;  
nothing is ever identical to anything else except itself.

— David LEWIS

On the Plurality of Worlds

• This is correct, but it is not the end of questions about identity

IDENTITY & {  
NECESSITY  
PROOF  
CONSTRUCTION  
ISOMORPHISM

} And philosophers have worked on  
these issues for a long time.

## IDENTITY & NECESSITY

8 = the number of planets ✓

It is necessary that  $8 = 8$  ✓

It is necessary that the number of planets = 8 ✗

$\Box[(\text{The } n \text{ where } n = \# \text{ planets}) n = 8]$  ✗ — de dicto

$(\text{The } n \text{ where } n = \# \text{ planets}) \Box[n = 8]$  ✓ — de re

Scope makes a difference



## IDENTITY & KNOWLEDGE / PROOF

Clark Kent = Superman. ✓

Lois Lane knows that Clark Kent is Clark Kent. ✓

Lois Lane knows that Clark Kent is Superman. ?

$f(x) = y$  ✓

Lois Lane shows that  $y = y$  ✓

Lois Lane shows that  $f(x) = y$  ?

$s = t$  's' & 't' have the same referent (value)  
they might not have the same sense.

# ISOMORPHISM & IDENTITY

Mathematical Structures — when are  $G_1$  &  $G_2$  the same group?

What is the relationship between isomorphism & identity?  
(this is a part of deciding what mathematical structure is.)

In philosophy of mathematics this is explored in  
**STRUCTURALISM**, which is congenial to category  
theoretical (& HoTT, Cubical) presentation, but  
these views are not identical.

(See, especially, Colin MacLarty, Steve Awodey.)

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# PROOF THEORY ♥ CONSTRUCTIVE LOGIC

Gentzen, Heyting, Dummett, PML, Prawitz, Girard

Understanding the Classical / Constructive boundary  
is an active research area in many directions.

Translations: Classical  $\overset{DN/\dots}{\hookrightarrow}$  Constructive  
Constructive  $\overset{S4}{\hookrightarrow}$  Classical Model  
<sub>Topological</sub>

The Context of Deduction:

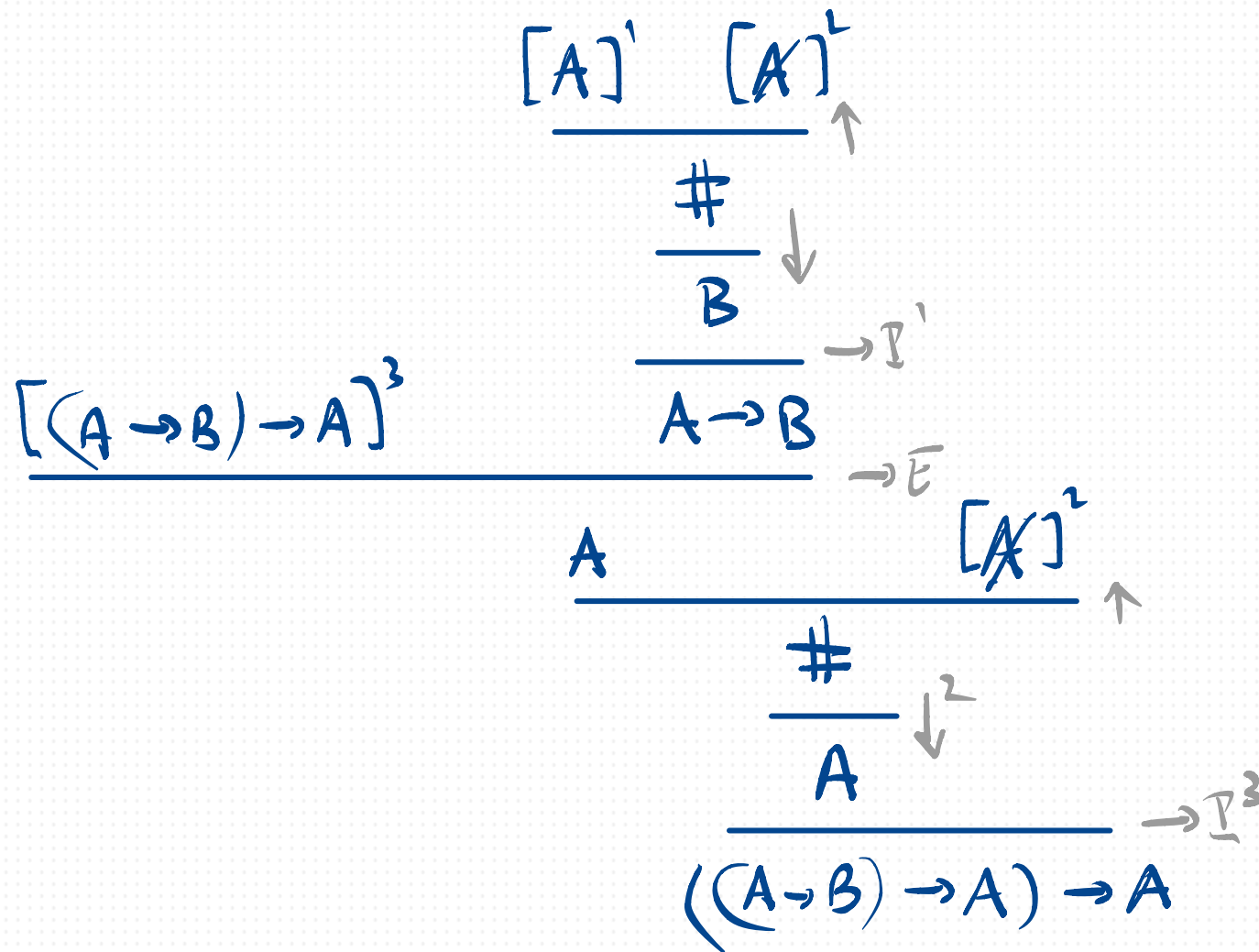
What is the difference between

$$\Gamma \vdash A \quad \&$$

$$\Gamma \vdash A, \Delta; \Gamma, A \vdash \Delta?$$

**BILATERALISM**: assertion & denial treated equally

# A (MILDLY) BILATERAL PROOF of PEIRCE'S LAW



Contexts contain positive & negative information. Judgements are positive or dead ends.

# A (MILDLY) BILATERAL PROOF of PEIRCE'S LAW

(In M-Parigot's  $\lambda\mu$  calculus)

$$\begin{array}{c}
 \frac{\kappa[A]^1 \quad \alpha[A]^1}{\phantom{\kappa[A]^1 \quad \alpha[A]^1}} \uparrow \\
 \frac{\alpha[x] \#}{\mu\beta.\alpha[x] \quad B} \downarrow \\
 \frac{\phantom{\kappa[A]^1 \quad \alpha[A]^1} \rightarrow \mathbb{I}'}{\phantom{\kappa[A]^1 \quad \alpha[A]^1} \rightarrow \mathbb{I}'} \\
 \frac{\gamma[(A \rightarrow B) \rightarrow A]^3 \quad \lambda x.\mu\beta.\alpha[x] \quad A \rightarrow B}{\phantom{\gamma[(A \rightarrow B) \rightarrow A]^3 \quad \lambda x.\mu\beta.\alpha[x] \quad A \rightarrow B}} \rightarrow \mathbb{E} \\
 \frac{\gamma(\lambda x.\mu\beta.\alpha[x]) \quad A \quad \alpha[A]^2}{\phantom{\gamma(\lambda x.\mu\beta.\alpha[x]) \quad A \quad \alpha[A]^2}} \uparrow \\
 \frac{\alpha[\gamma(\lambda x.\mu\beta.\alpha[x])] \#}{\mu\alpha.\alpha[\gamma(\lambda x.\mu\beta.\alpha[x])] \quad A} \downarrow^2 \\
 \phantom{\mu\alpha.\alpha[\gamma(\lambda x.\mu\beta.\alpha[x])] \quad A} \rightarrow \mathbb{I}^3 \\
 \lambda y.\mu\alpha.\alpha[\gamma(\lambda x.\mu\beta.\alpha[x])] \quad ((A \rightarrow B) \rightarrow A) \rightarrow A
 \end{array}$$

Contexts contain positive & negative information. Judgements are positive or dead ends.

# A SYMMETRIC BILATERAL CALCULUS $\bar{\lambda}\mu\tilde{\mu}$ (CURIE & HERBELIN)

At the typing level, we obtain  $LK_{\mu\tilde{\mu}}$  whose typing judgements are:

$c : (\Gamma \vdash \Delta)$	COMMANDS		CLASHES
$\Gamma \vdash v : A \mid \Delta$	TERMS		ASSERTIONS
$\Gamma \mid e : A \vdash \Delta$	CONTEXTS		DENIALS

and whose typing rules are:

$$\frac{\Gamma \vdash v : A \mid \Delta \quad \Gamma \mid e : A \vdash \Delta}{\langle v \mid e \rangle : (\Gamma \vdash \Delta)}$$

$$\frac{}{\Gamma \mid \alpha : A \vdash \alpha : A, \Delta}$$

$$\frac{}{\Gamma, x : A \vdash x : A \mid \Delta}$$

$$c : (\Gamma \vdash \beta : B, \Delta)$$

$$\Gamma \vdash \mu\beta.c : B \mid \Delta$$

$$c : (\Gamma, x : A \vdash \Delta)$$

$$\Gamma \mid \tilde{\mu}x.c : A \vdash \Delta$$

$$\frac{\Gamma \vdash v : A \mid \Delta \quad \Gamma \mid e : B \vdash \Delta}{\Gamma \mid v \cdot e : A \rightarrow B \vdash \Delta}$$

$$\Gamma, x : A \vdash v : B \mid \Delta$$

$$\Gamma \vdash \lambda x.v : A \rightarrow B \mid \Delta$$

complex contexts / denials.

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## FREGE'S BEGRIFFSCHRIFT

$A$  — content, a thought, a proposition.

$\vdash A$  — the assertion that  $A$

If  $A$  then  $B$ . — this can be asserted, but the  $A$  &  $B$  are propositions inside the conditional, but are not asserted.

Assertion is a speech act — there are others.

$?A$  — polar question

$?_x A(x)$  — find an  $x$  where  $A(x)$  question

$!A$  — see to it that  $A$  is true

$\{_\beta A$  —  $\beta$  promises to see to it that  $A$

} QUESTION

} IMPERATIVE

} COMMITMENT

## CONDITIONAL SPEECH ACTS

If A then is it the case that B?

If A then I promise to B.

If A then please do B.

Are these questions,  
promises & imperatives?

Certainly if the antecedent  
holds... maybe only then.

If A is a restrictor of more than propositions.

Traditional formal grammars do not respect conditional  
speech acts — the grammar is independent of the semantics.

[A true]		[A true]	
⋮		⋮	
A prop	B prop	A prop	B promise
<hr/>		<hr/>	
A > B prop		A > B promise	

These are also entangled, but the dependence is in the other direction!

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# TYPE THEORY CAN BE APPLIED IN DIFFERENT WAYS

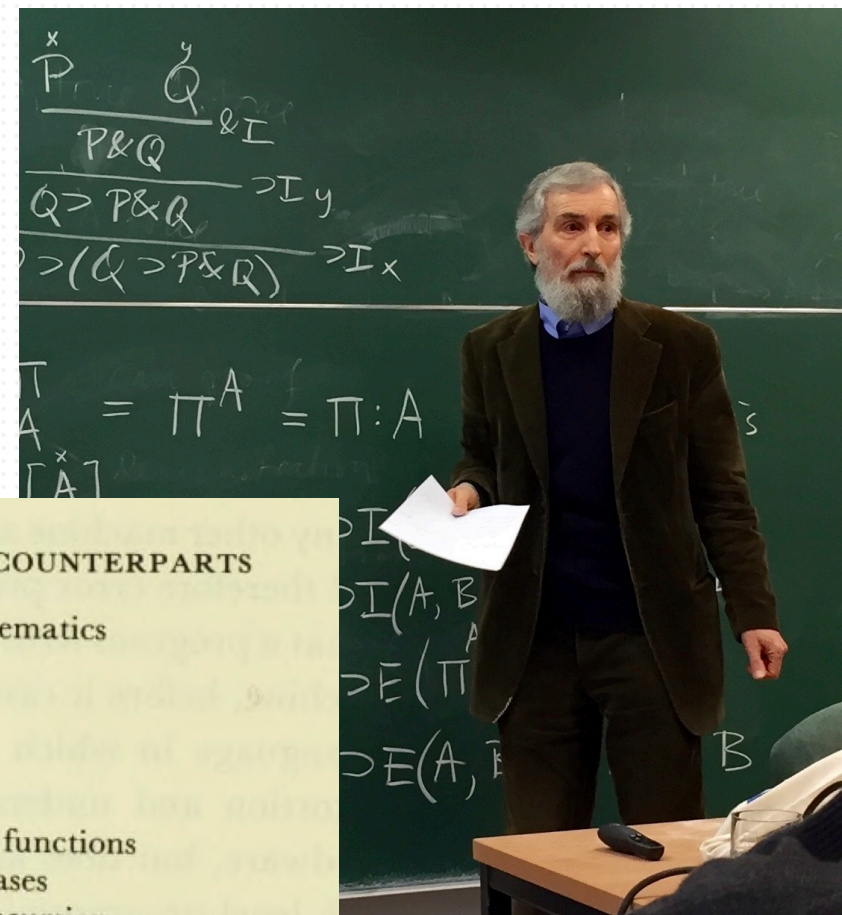


TABLE 1. KEY NOTIONS OF PROGRAMMING WITH MATHEMATICAL COUNTERPARTS

programming	mathematics
program, procedure, algorithm	function
input	argument
output, result	value
$x := e$	$x = e$
$S_1; S_2$	composition of functions
<b>if</b> $B$ <b>then</b> $S_1$ <b>else</b> $S_2$	definition by cases
<b>while</b> $B$ <b>do</b> $S$	definition by recursion
data structure	element, object
data type	set, type
value of a data type	element of a set, object of a type
$a:A$	$a \in A$
<b>integer</b>	$\mathbb{Z}$
<b>real</b>	$\mathbb{R}$
<b>Boolean</b>	$\{0, 1\}$
$(c_1, \dots, c_n)$	$\{c_1, \dots, c_n\}$
<b>array</b> $[I]$ <b>of</b> $T$	$T^I, I \rightarrow T$
<b>record</b> $s_1:T_1; s_2:T_2$ <b>end</b>	$T_1 \times T_2$
<b>record case</b> $s:(c_1, c_2)$ <b>of</b> $c_1:(s_1:T_1); c_2:(s_2:T_2)$ <b>end</b>	$T_1 + T_2$
<b>set of</b> $T$	$\{0, 1\}^T, T \rightarrow \{0, 1\}$

PMc - Constructive Mathematics & Computer programming (1984)

Computational type theory  
sequents classify computational processes

Formal type theory  
pure logic, backed only by the rules

$$\Gamma \vdash t : A,$$

...

Dialogical type theory  
sequents classify practices of  
processes of reasoning & justification

NORMATIVE PRAGMATICS

Robert Brandom, Jaroslav Peregrin, ...

Conceptual type theory  
sequents classify cognitive constructions

CONSTRUCTIVE LOGIC

INTUITIONISM — theories of judgement

PML, Dag Prawitz, Göran Sundholm, ...

# HYBRID TYPE THEORY?

Computational type theory  
sequents classify computational processes

Formal type theory

$$\Gamma \vdash t : A$$

...

What about applications that encompass these domains? Justifications that include computation, computer aided reasoning, natural language program specification

Dialogical type theory  
sequents classify practices & processes of reasoning & justification

Conceptual type theory  
sequents classify cognitive constructions

It seems to me that many of these intersections could be fruitful in the years ahead.

Questions?

