

About the construction of simplicial and cubical sets in indexed form: the case of degeneracies

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The problem of *semi-simplicial types*

Construct the following family of family types in homotopy type theory:

$X_0 : \mathbf{Type}$ (points)

$X_1 : X_0 \times X_0 \rightarrow \mathbf{Type}$ (segments)

$X_2 : \prod x_1 x_2 x_3. X_1(x_1, x_2) \rightarrow X_1(x_1, x_3) \rightarrow X_1(x_2, x_3) \rightarrow \mathbf{Type}$ (triangles)

\vdots

The problem of *semi-simplicial types*

Construct the following family of family types in pure type theory:

$X_0 : \mathbf{Type}$ (points)

$X_1 : X_0 \times X_0 \rightarrow \mathbf{Type}$ (segments)

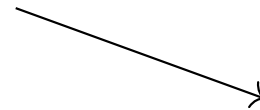
$X_2 : \prod x_1 x_2 x_3. X_1(x_1, x_2) \rightarrow X_1(x_1, x_3) \rightarrow X_1(x_2, x_3) \rightarrow \mathbf{Type}$ (triangles)

\vdots

two problems into one



semi-simplicial *types*
needs – at least – a
description of higher-
dimensional coherences



recipe for *indexed* semi-simplicial
this talk, for **HSet**,
see also Voevodsky, Part-Luo,
Altenkirch-Capriotti-Kraus

The fibred/indexed correspondence for **HSet**

For $B : \mathbf{HSet}$

$$\begin{array}{ccc} E : \mathbf{HSet} & & \\ \downarrow & \simeq & B \rightarrow \mathbf{HSet} \\ B & & \end{array}$$

Iterating the fibred/indexed correspondence for **HSet**

Application to definition of Reedy presheaves in *indexed form*, here for semi-cubical sets:

$\begin{array}{c} \text{fibred form} \\ Y_0 : \mathbf{HSet} \end{array}$	vs	$\begin{array}{c} \text{indexed form} \\ X_0 : \mathbf{HSet} \end{array}$	
$\begin{array}{c} \uparrow d^L \uparrow d^R \\ Y_1 : \mathbf{HSet} \end{array}$		$X_1 : X_0 \times X_0 \rightarrow \mathbf{HSet}$	
$\begin{array}{c} \uparrow d^{L*} \uparrow d^{R*} \uparrow d^{*L} \uparrow d^{*R} \\ Y_2 : \mathbf{HSet} \\ + \text{coherences} \end{array}$		$\begin{array}{l} X_2 : \Pi(x_{LL}, x_{LR}). \Pi x_{L*} : X_1(x_{LL}, x_{LR}). \\ \Pi(x_{RL}, x_{RR}). \Pi x_{R*} : X_1(x_{RL}, x_{RR}). \\ X_1(x_{LL}, x_{RL}) \times X_1(x_{LR}, x_{RR}) \rightarrow \mathbf{HSet} \end{array}$	

Iterating the fibred/indexed correspondence for **HSet**

Application to the definition of Reedy presheaves in *indexed form*, here for cubical sets:

<i>fibred form</i>	<i>vs</i>	<i>indexed form</i>
$Y_0 : \mathbf{HSet}$		$X_0 : \mathbf{HSet}$
$\begin{array}{c} \uparrow \\ d^L \\ \uparrow \end{array} \quad \begin{array}{c} \uparrow \\ d^R \\ \uparrow \end{array}$		
$Y_1 : \mathbf{HSet}$		$X_1 : X_0 \times X_0 \rightarrow \mathbf{HSet}$
$\begin{array}{c} \uparrow \\ d^{L*} \\ \uparrow \end{array} \quad \begin{array}{c} \uparrow \\ d^{R*} \\ \uparrow \end{array} \quad \begin{array}{c} \uparrow \\ d^{\star L} \\ \uparrow \end{array} \quad \begin{array}{c} \uparrow \\ d^{\star R} \\ \uparrow \end{array}$		
$Y_2 : \mathbf{HSet}$		$X_2 : \Pi(x_{LL}, x_{LR}). \Pi x_{L*} : X_1(x_{LL}, x_{LR}).$ $\Pi(x_{RL}, x_{RR}). \Pi x_{R*} : X_1(x_{RL}, x_{RR}).$ $X_1(x_{LL}, x_{RL}) \times X_1(x_{LR}, x_{RR}) \rightarrow \mathbf{HSet}$
+ coherences		

Motivations:

1. The iterated fibred/indexed correspondence is interesting in itself
2. Suggests models of type theory closer to the syntax: e.g. equality interpreted as a (relevant) relation rather than as a span

Rest of the talk

Presheaves in “indexed” form

- following a n -ary “parametricity” recipe
 - s.t. *unary* parametricity gives augmented semi-simplicial sets
 - and *binary* parametricity gives semi-cubical sets
- equipped with a *degeneracy*
- machine-checked in Rocq

A uniform approach to augmented simplicial sets and cubical sets

Augmented simplicial and cubical categories only differ in the “arity” of a finite set ν :

$$\mathbf{Obj} \quad := \mathbb{N}$$

$$\mathbf{Hom}(p, n) := \{l \in (\nu \sqcup \{\star\})^n \mid \text{number of } \star \text{ in } l = p\}$$

$$g \circ f \quad := \begin{cases} f & \text{if } g = \epsilon \\ a(g' \circ f) & \text{if } g = a g', \text{ where } a \in \nu \\ a(g' \circ f') & \text{if } g = \star g', f = a f', \text{ where } a \in \nu \text{ or } a = \star \end{cases}$$

$$\mathbf{id} \quad := \star \dots \star \quad n \text{ times for } \mathbf{id} \in \mathbf{Hom}(n, n)$$

A uniform approach to augmented simplicial sets and cubical sets

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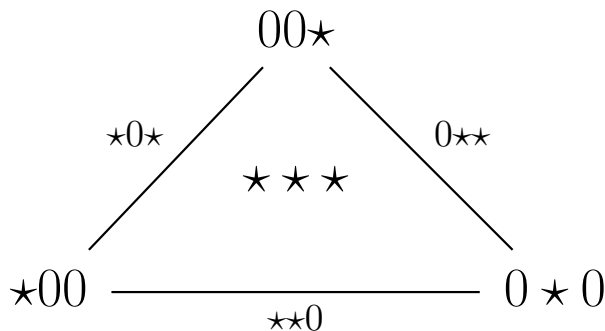
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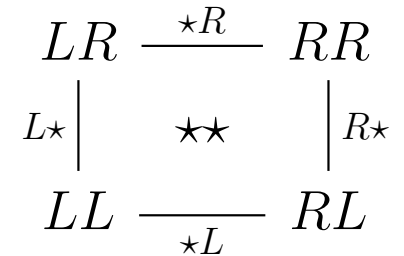
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$$\text{id} \quad := \star \dots \star \quad n \text{ times for } \text{id} \in \text{Hom}(n, n)$$

augmented semi-simplicial sets with $\nu = \{0\}$
(counting from -1)



semi-cubical sets with $\nu = \{L, R\}$
(counting from 0)



An effective indexed construction as a dependent stream of dependent sets

ν -sets

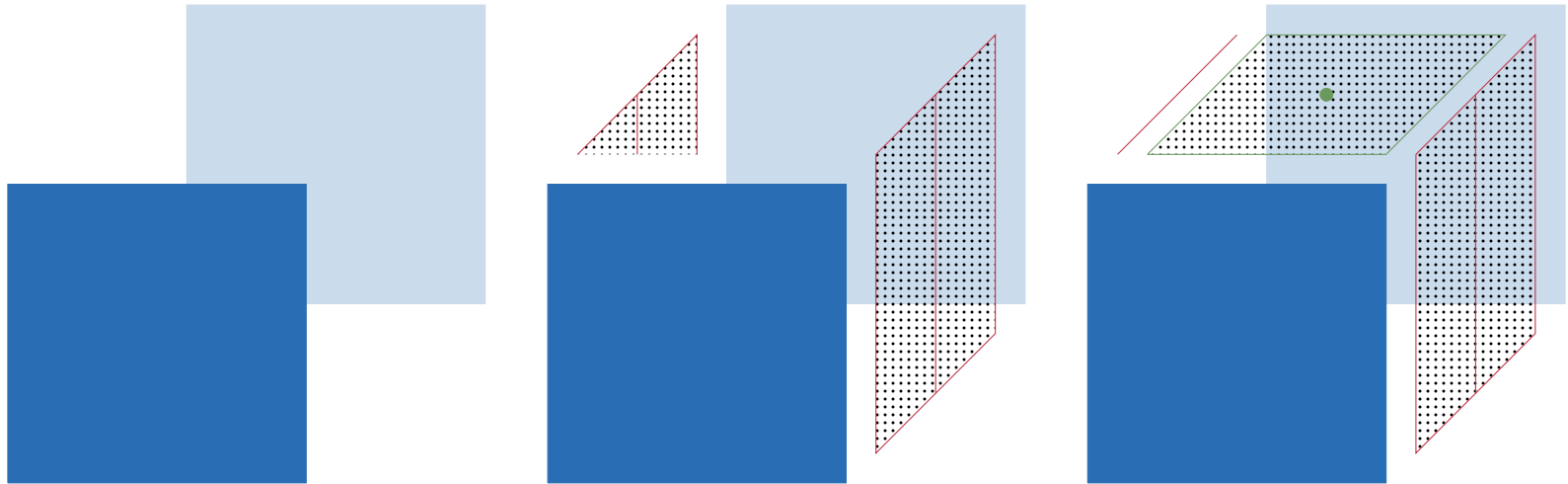
νSet_m	:	HSet_{m+1}
νSet_m	\triangleq	$\nu\text{Set}_m^{\geq 0}(\ast)$
$\nu\text{Set}_m^{\geq n} \quad (D : \nu\text{Set}_m^{< n})$:	HSet_{m+1}
$\nu\text{Set}_m^{\geq n} \quad D$	\triangleq	$\Sigma R : \nu\text{Set}_m^{=n}(D). \nu\text{Set}_m^{\geq n+1}(D, R)$

Truncated ν -sets

$\nu\text{Set}_m^{< n}$:	HSet_{m+1}
$\nu\text{Set}_m^{< 0}$	\triangleq	unit
$\nu\text{Set}_m^{< n'+1}$	\triangleq	$\Sigma D : \nu\text{Set}_m^{< n'} . \nu\text{Set}_m^{=n'}(D)$
$\nu\text{Set}_m^{=n} \quad (D : \nu\text{Set}_m^{< n})$:	HSet_{m+1}
$\nu\text{Set}_m^{=n} \quad D$	\triangleq	fullframe $_m^n(D) \rightarrow \text{HSet}_m$

where **fullframe** $_m^n$ is defined by mutual recursive construction (see next slides)

The recursive process used to build frames from layers of paintings



The recursive construction, formally

fullframe^n	$(D : \text{vSet}_m^{<n})$:	HSet_m
fullframe^n	D	\triangleq	$\text{frame}^{n,n}(D)$
$\text{frame}^{n,p,p \leq n}$	$(D : \text{vSet}_m^{<n})$:	HSet_m
$\text{frame}^{n,0}$	D	\triangleq	unit
$\text{frame}^{n,p'+1}$	D	\triangleq	$\Sigma d : \text{frame}^{n,p'}(D). \text{layer}^{n,p'}(d)$
$\text{layer}^{n,p,p < n}$	$\{D : \text{vSet}_m^{<n}\}$ $(d : \text{frame}^{n,p}(D))$:	HSet_m
$\text{layer}^{n,p}$	$D \ d$	\triangleq	$\Pi \omega. \text{painting}^{n-1,p}(D.2)(\text{restr}_{\text{frame},\omega,p}^{n,p}(d))$
$\text{painting}^{n,p,p \leq n}$	$(D : \text{vSet}_m^{<n})$ $(E : \text{vSet}_m^{=n}(D))$ $(d : \text{frame}^{n,p}(D))$:	HSet_m
$\text{painting}^{n,p,p=n}$	$D \ E \ d$	\triangleq	$E(d)$
$\text{painting}^{n,p,p < n}$	$D \ E \ d$	\triangleq	$\Sigma l : \text{layer}^{n,p}(d). \text{painting}^{n,p+1}(E)(d, l)$

where we need to define $\text{restr}_{\text{frame}}$ (see next slide)

The recursive construction: restrictions (“faces”)

$\text{restr}_{\text{frame}, e, q}^{n, p, p \leq q \leq n-1}$	$\{D : \text{vSet}^{<n}\}$ $(d : \text{frame}^{n, p}(D))$:	$\text{frame}^{n-1, p}(D.1)$
$\text{restr}_{\text{frame}, e, q}^{n, 0}$	$D *$	\triangleq	$*$
$\text{restr}_{\text{frame}, e, q}^{n, p'+1}$	$D (d, l)$	\triangleq	$(\text{restr}_{\text{frame}, e, q}^{n, p'}(d), \text{restr}_{\text{layer}, e, q-1}^{n, p'}(l))$
$\text{restr}_{\text{layer}, e, q}^{n, p, p \leq q \leq n-2}$	$\{D : \text{vSet}^{<n}\}$ $\{d : \text{frame}^{n, p}(D)\}$ $(l : \text{layer}^{n, p}(d))$:	$\text{layer}^{n-1, p}(\text{restr}_{\text{frame}, e, q+1}^{n, p}(d))$
$\text{restr}_{\text{layer}, e, q}^{n, p}$	$D d l$	\triangleq	$\lambda \omega. (\text{coh}_{\text{frame}, e, \omega, q, p}^{n, p}(d) (\text{restr}_{\text{painting}, e, q}^{n-1, p}(D.2)(l_\omega)))$
$\text{restr}_{\text{painting}, e, q}^{n, p, p \leq q \leq n-1}$	$(D : \text{vSet}^{<n})$ $(E : \text{vSet}^{=n}(D))$ $(d : \text{frame}^{n, p}(D))$ $(c : \text{painting}^{n, p}(E)(d))$:	$\text{painting}^{n-1, p}(D.2)(\text{restr}_{\text{frame}, e, q}^{n, p}(d))$
$\text{restr}_{\text{painting}, e, q}^{n, p, p=q}$	$D E d (l, _)$	\triangleq	l_e
$\text{restr}_{\text{painting}, e, q}^{n, p, p < q}$	$D E d (l, c)$	\triangleq	$(\text{restr}_{\text{layer}, e, q-1}^{n, p}(l), \text{restr}_{\text{painting}, e, q}^{n, p+1}(E)(c))$

where we need to define $\text{coh}_{\text{frame}}$ (see next slide)

The recursive construction: coherences

$\text{coh}_{\text{frame}, \varepsilon, \omega, q, r}^{n, p, p \leq r \leq q \leq n-2}$	$\{D : \nu\text{Set}^{<n}\}$ $(d : \text{frame}(D))$:	$\text{restr}_{\text{frame}, \varepsilon, q}^{n-1, p}(\text{restr}_{\text{frame}, \omega, r}^{n, p}(d))$ $= \text{restr}_{\text{frame}, \omega, r}^{n-1, p}(\text{restr}_{\text{frame}, \varepsilon, q+1}^{n, p}(d))$
$\text{coh}_{\text{frame}, \varepsilon, \omega, q, r}^{n, 0}$	$D *$	\triangleq	$\text{refl}(*)$
$\text{coh}_{\text{frame}, \varepsilon, \omega, q, r}^{n, p'+1}$	$D(d, l)$	\triangleq	$(\text{coh}_{\text{frame}, \varepsilon, \omega, q, r}^{n, p'}(d), \text{coh}_{\text{layer}, \varepsilon, \omega, q-1, r-1}^{n, p'}(l))$
$\text{coh}_{\text{layer}, \varepsilon, \omega, q, r}^{n, p, p \leq r \leq q \leq n-3}$	$\{D : \nu\text{Set}^{<n}\}$ $\{d : \text{frame}(D)\}$ $(l : \text{layer}(d))$:	$\text{restr}_{\text{layer}, \varepsilon, q}^{n-1, p}(\text{restr}_{\text{layer}, \omega, r}^{n, p}(l))$ $= \text{restr}_{\text{layer}, \omega, r}^{n-1, p}(\text{restr}_{\text{layer}, \varepsilon, q+1}^{n, p}(l))$
$\text{coh}_{\text{layer}, \varepsilon, \omega, q, r}^{n, p}$	$D d l$	\triangleq	$\lambda \theta. \text{coh}_{\text{painting}, \varepsilon, \omega, q, r}^{n-1, p}(D.2)(l_\theta)$
$\text{coh}_{\text{painting}, \varepsilon, \omega, q, r}^{n, p, p \leq r \leq q \leq n-2}$	$\{D : \nu\text{Set}^{<n}\}$ $(E : \nu\text{Set}^{=n}(D))$ $\{d : \text{frame}(D)\}$ $(c : \text{painting}(E)(d))$:	$\text{restr}_{\text{painting}, \varepsilon, q}^{n-1, p}(D.2)(\text{restr}_{\text{painting}, \omega, r}^{n, p}(E)(c))$ $= \text{restr}_{\text{painting}, \omega, r}^{n-1, p}(D.2)(\text{restr}_{\text{painting}, \varepsilon, q+1}^{n, p}(E)(c))$
$\text{coh}_{\text{painting}, \varepsilon, \omega, q, r}^{n, p, p=r}$	$D E d(l, _)$	\triangleq	$\text{refl}(\text{restr}_{\text{painting}, \varepsilon, q}^{n-1, p}(D.2)(l_\omega))$
$\text{coh}_{\text{painting}, \varepsilon, \omega, q, r}^{n, p, p < r}$	$D E d(l, c)$	\triangleq	$(\text{coh}_{\text{layer}, \varepsilon, \omega, q-1, r-1}^{n, p}(l), \text{coh}_{\text{painting}, \varepsilon, \omega, q, r}^{n, p+1}(E)(c))$

where we hide many steps of equational reasoning: proof-irrelevance of equality in \mathbf{HSet} , identification of equality of pairs and pairs of equalities, groupoid properties of equality

About the formalisation

Complex proof of termination

- made several unsuccessful attempts
- construction completed in Rocq in Apr 2022
(inductively building 3 levels at once with two subinductions)
- degeneracies completed in Nov 2024
- we are working on a simplification saving a lot of equational reasoning
- code at <https://github.com/artagnon/bonak>

Note: “paper” construction also fully formulated in Agda (w/o termination)

Adding (one) degeneracy (in the last direction)

fibred form *vs* *indexed form*

Y_0 : HSet	X_0 : HSet	(points)
$\uparrow\uparrow\downarrow$		
Y_1 : HSet	$X_1 : X_0 \times X_0 \rightarrow \text{HSet}$	(segments)
$\uparrow\uparrow\uparrow\uparrow\downarrow$	$r_0 : \prod x_0 : X_0. X_1(x_0, x_0)$	
Y_2 : HSet	$X_2 : \prod(x_{LL}^0, x_{LR}^0). \prod x_{L*}^1 : X_1(x_{LL}^0, x_{LR}^0).$ $\prod(x_{RL}^0, x_{RR}^0). \prod x_{R*}^1 : X_1(x_{RL}^0, x_{RR}^0).$ $X_1(x_{LL}^0, x_{RL}^0) \times X_1(x_{LR}^0, x_{RR}^0) \rightarrow \text{HSet}$	(squares)
+ coherences	$r_1 : \prod(x_L^0, x_R^0) : (X_0 \times X_0). \prod x^1 : X_1(x_L^0, x_R^0).$ $X_2((x_L^0, x_L^0), r_0(x_L^0), (x_R^0, x_R^0), r_0(x_R^0), (x^1, x^1))$	
\vdots	\vdots	

The added degeneracy is *parametric*: in the binary case, it gives a standard cubical degeneracy; in the unary case, it gives a ParamTT-like degeneracy and *not* a simplicial degeneracy

First, our degeneracy implies a distinguished point $r_{-1}(a)$ for any $a : X_{-1}$. Then:

source

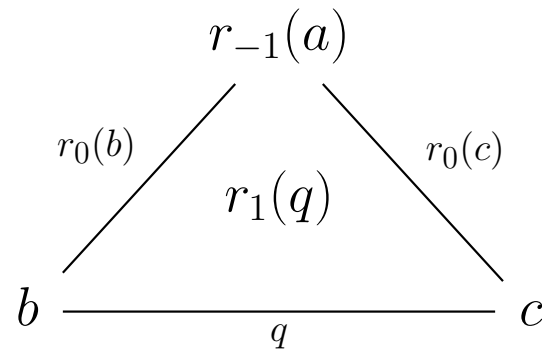
(over some $a : X_{-1}$)

b

$b \xrightarrow{q} c$

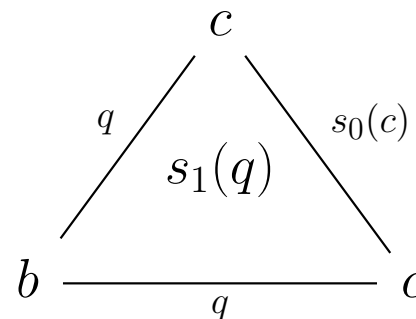
parametric
degeneracy

$b \xrightarrow{r_0(b)} r_{-1}(a)$



simplicial
degeneracy

$b \xrightarrow{s_0(b)} b$



actually
a 1-connection!

Adding a degeneracy

For any $(X_0, X_1, \dots) : \nu\mathbf{Set}$, we define a stream of degeneracies:

$$\begin{aligned} \nu\mathbf{reflSet}(X_0, X_1, \dots) &\triangleq \\ \Sigma r_0 : \Pi d : \mathbf{fullframe}^0. \Pi x : X_0(d). X_1(\mathbf{refl}_{\mathbf{fullframe}}^0(d), \lambda\epsilon. x). \\ \Sigma r_1 : \Pi d : \mathbf{fullframe}^1(X_0). \Pi x : X_1(d). X_2(\mathbf{refl}_{\mathbf{fullframe}}^1(r_0)(d), \lambda\epsilon. x). \\ \Sigma r_2 : \Pi d : \mathbf{fullframe}^2(X_0, X_1). \Pi x : X_2(d). X_3(\mathbf{refl}_{\mathbf{fullframe}}^2(r_0, r_1)(d), \lambda\epsilon. x). \\ \dots \end{aligned}$$

where

$$\mathbf{refl}_{\mathbf{fullframe}}^n(r_{-1}, \dots, r_{n-1}) : \mathbf{fullframe}^n(X_{-1}, \dots, X_{n-1}) \rightarrow \mathbf{frame}^{n+1, n}(X_{-1}, \dots, X_n)$$

computes the n first layers of the border of $r_n(d)(x)$, knowing that the last layer is made of ν times x itself, so that

$$(\mathbf{refl}_{\mathbf{fullframe}}^n(r_{-1}, \dots, r_{n-1})(d), \lambda\epsilon. x) : \mathbf{frame}^{n+1, n+1}(X_{-1}, \dots, X_n)$$

is a full frame.

Adding a degeneracy

On the way, we need two coherence conditions:

$$\text{idrestrrefl}_{\text{frame}, \epsilon}^n(r_{-1}, \dots, r_{n-1}) (d : \text{fullframe}^n(X_0, \dots, X_{n-1})) : \\ \text{restr}_{\text{frame}, \epsilon, n}^{n, n}(\text{refl}_{\text{fullframe}}^n(r_{-1}, \dots, r_{n-1})(d)) = d$$

$$\text{cohrestrrefl}_{\text{frame}, \epsilon, p < n}^n(r_{-1}, \dots, r_{n-1}) (d : \text{frame}^{n, p}(X_0, \dots, X_{n-1})) : \\ \text{restr}_{\text{frame}, \epsilon, p}^{n, p}(\text{refl}_{\text{frame}}^{n, p}(r_{-1}, \dots, r_{n-1})(d)) = \text{refl}_{\text{frame}}^{n-1, p}(r_{-1}, \dots, r_{n-2})(\text{restr}_{\text{frame}, \epsilon, p}^{n-1, p}(d))$$

where $\text{refl}_{\text{frame}}^{n, p}$ generalises $\text{refl}_{\text{fullframe}}^n$ to prefixes of fullframe^n :

$$\text{refl}_{\text{frame}}^{n, p}(r_{-1}, \dots, r_{n-1}) : \text{frame}^{n, p}(X_{-1}, \dots, X_{n-1}) \rightarrow \text{frame}^{n+1, p}(X_{-1}, \dots, X_n)$$

Summary

- Machine-checked parametricity-based definition of *indexed* presheaves
- Uniformly represents simplicial and cubical sets
- Addition of one (parametric) degeneracy in the last direction completed
- More compact definition in progress, relying on finer-grain dependencies between the different components of the construction